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# Self-similarity in Parallel I/Os

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# I/O Arrivals in Scientific Applications

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- Understanding I/O behavior of parallel scientific applications is important for tuning, managing, or optimizing parallel file systems;
- F. Wang et al (Ref. [8]) examine the I/O burstiness of parallel I/O workloads using a simple methodology. They measure the cumulative distribution functions (CDF) of I/O inter-arrival times and conclude that *I/O activities in the LLNL traces are very bursty in the ior2 benchmark and the f1 application.*

# Motivation

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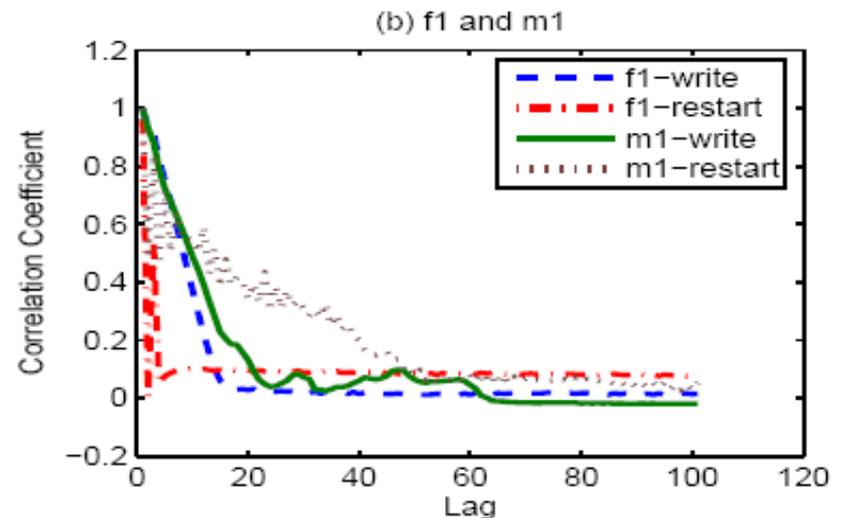
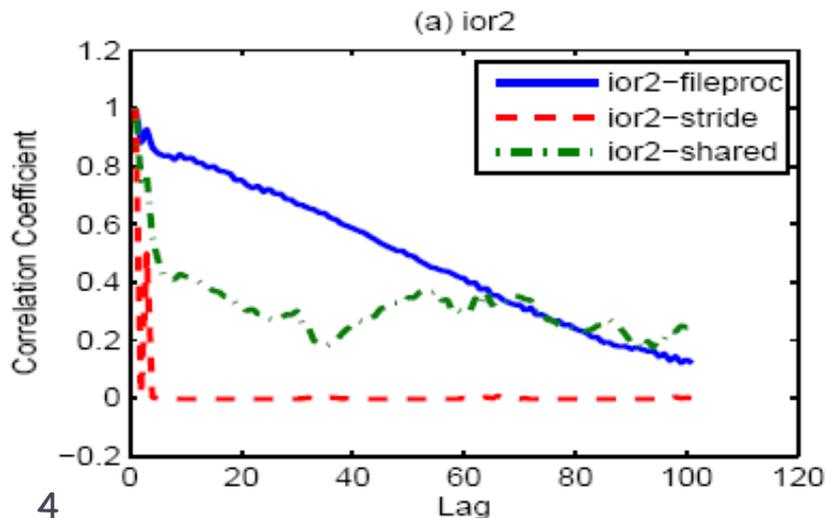
- Is it appropriate to use a Poisson or Markov model to characterize or predict parallel I/O arrivals with the presence of intensive burstiness in scientific applications?
  - ✓ In order to model parallel I/O workloads, prior studies usually assumed that *I/O arrival process follows a Poisson distribution*, and I/Os can be generated by using the Markov Model;
  - ✓ However, the Markov model does not accurately characterize the *burstiness* in parallel I/O workloads;
  - ✓ There are very bursty I/O activities as evidenced in the LLNL application workloads, such as the *ior2* benchmark and the *f1* application described in Ref. [8].

# Auto-Correlation of Inter-arrival Times

- We use *auto-correlation functions* (ACF) to measure if earlier values in a time sequence have some correlation to later values

$$ACF(k) = \frac{c_k}{c_0} \quad \text{where} \quad c_k = \frac{1}{N-k} \sum_{i=1}^{N-k} (x_i - \bar{x})(x_{i+k} - \bar{x})$$

- It might be appropriate to use Markov model to synthesize *ior2-stride* workload, but not *ior2-fileproc*, *ior2-shared*;
- Examination results above motivate us to further study the self-similarity of parallel I/Os.



# Self-similarity Detection

➤ A covariance stationary stochastic process is self-similar if

$$\lim_{k \rightarrow \infty} \frac{ACF(k)}{k^{-\beta}} = c < \infty, \text{ for } 0 < \beta < 1$$

➤ Second order self-similar  $ACF(k) = \frac{1}{2} [(k+1)^{2-\beta} - 2k^{2-\beta} + (k-1)^{2-\beta}]$

➤ Hurst parameter,  $H = 1 - \beta/2$ , measures of the degree of self-similarity.

➤ Use three methods to estimate H: variance-time plot (VTP), R/S, and Whittle estimator

➤ Most of the Hurst exponent values are above 0.5 except the I/O events in ior2-stride. This observation indicates the comprehensive *existence of self-similarity* in the LLNL I/O traces studied in this paper.

Traces	Estimation of $H$				
Streams	1	2	3	4	5
<i>ior2-fileproc</i>	0.67	0.66	0.65	0.69	0.70
<i>ior2-shared</i>	0.72	0.78	0.83	0.84	0.86
<i>ior2-stride</i>	0.52	0.49	0.45	0.40	0.41
<i>fl-write</i>	0.55	0.56	0.54	0.58	0.56
<i>fl-restart</i>	0.51	0.59	0.52	0.50	0.64
<i>m1-write</i>	0.66	0.67	0.67	0.65	0.66
<i>m1-restart</i>	0.67	0.65	0.66	0.64	0.63

# Modeling I/O Arrivals

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➤ We find that the parallel I/O traces studied meet the major requirements of a widely used model named *Fractional Brownian Motion*.

➤ The prediction algorithm can be expressed as

$$\lambda^{2H-1} \sim \frac{-2\alpha \ln(\epsilon) ((1-H)^{1-H} H^H)^2 b^{2H-2}}{[b(1-\epsilon) - \epsilon]^{2H}}.$$

➤ Our prediction model can be briefly described as right.

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## ARRIVAL-RATE-PREDICTION

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**INPUT:** The I/O miss probability  $\epsilon$ , the maximum number of outstanding streams served at per node  $b$ , original trace data file  $f$ .

**OUTPUT:** arrival rate  $\{\lambda(i); i = 1, 2, \dots, n\}$ .

**ALGORITHM:**

**for** each  $f$

Use maximum-likelihood estimate to estimate the parameter value  $\alpha$  for data sets in  $f$ ;  
Use Pox plot to estimate the Hurst value  $H$

**if**  $H \notin (0, 1)$  or  $H = 1/\alpha$

**then** break;

**else**

Set the initial values of  $\epsilon$  and  $b$ , and obtain  $\{\lambda(i); i = 1, 2, \dots, n\}$  using Equation (5)

**end for**

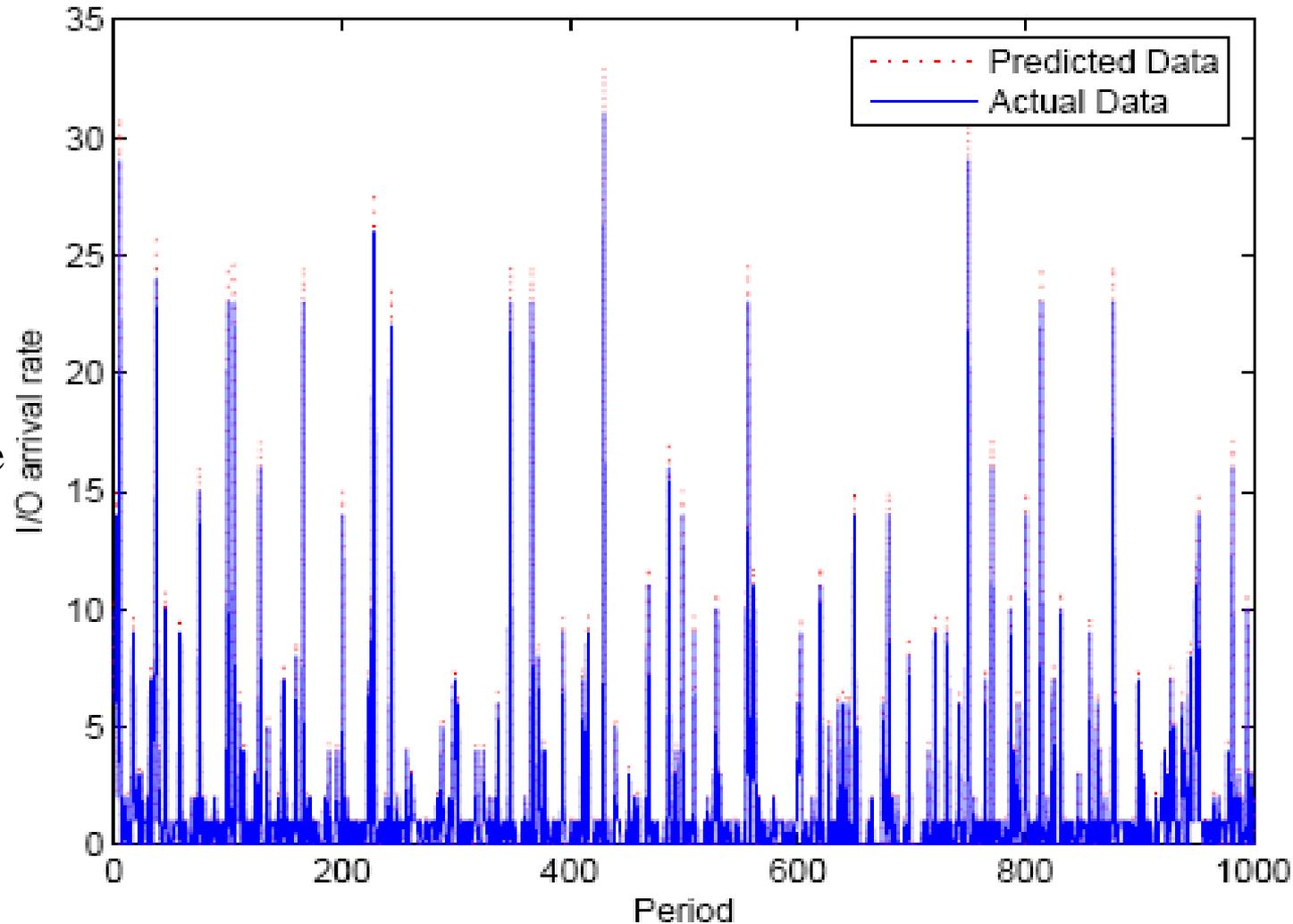
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# Evaluation of Our Model

➤ Our model vs realistic LLNL I/O traces

➤ Results of *one randomly selected stream* in the figure right.

➤ Both sequences have exactly the same mean arrival rate.



# Conclusions

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- There are *evident correlations between inter-arrival times* in subtraces collected on most computing nodes.
- Scientific I/O workloads are self-similar for short-term scales. Thus traditional *Poisson or Markovian arrival processes are inappropriate* to model the I/O demands.
- We *develop an accurate analytical model* to model I/O mean arrival rate for I/O workloads with self-similarity.
- Our immediate future work is to collect parallel I/O traces lasting weeks or months and further evaluate self-similarity.