

Parallel Reed/Solomon Coding on Multicore Processors

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Erasure-tolerant Codes

k data storage resources, e.g. disks

m redundant resources

regular data striping across k resources

encoding: calculation of m independent redundant blocks



a code tolerates f failed storage resources: $f \leq m$

Criteria:

- Number tolerated faults: $f = m$ as the optimum
- Flexibility, when choosing k , m
- Computational cost for en- and decoding

Cauchy Reed/Solomon

Encoding:

- Multiplication of original data word o with a generator matrix G

$$a = \begin{bmatrix} o \\ c \end{bmatrix} = G \cdot o = \begin{bmatrix} I \\ G_{sub} \end{bmatrix} \cdot o$$

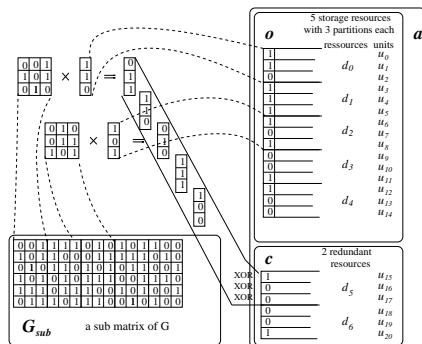
Example Reed/Solomon, 5+2:

- operations $+$, \cdot within $GF(2^3)$

$$c = \left\{ \begin{array}{ccccc} 2 & 7 & 4 & 3 & 1 \\ 3 & 4 & 7 & 2 & 5 \end{array} \right\} \cdot o$$

as a Cauchy-Reed/Solomon code:

- projection to $GF(2^1)$, binary logic
- operations XOR, AND



Cauchy Reed/Solomon

Decoding:

- equations system used for data recalculation

$$o = G'^{-1} * a'$$

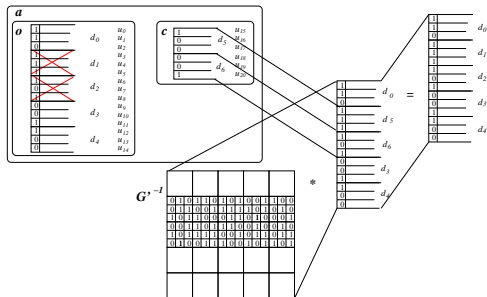
when 2nd and 3rd resource fail:

- operations $+$, \cdot within $GF(2^3)$

$$o = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 4 & 5 & 3 & 2 & 1 \\ 2 & 3 & 5 & 4 & 4 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} d0 \\ d5 \\ d6 \\ d3 \\ d4 \end{pmatrix}$$

by Cauchy-Reed/Solomon:

- projection to $GF(2^1)$, binary logic
- operations XOR, AND



Cauchy Reed/Solomon: Equation-based definition

Equations refer different bits within storage resources

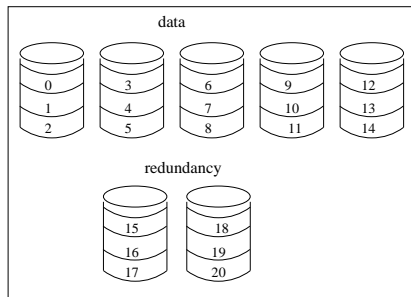
different bits \Rightarrow units on resources \Rightarrow partitions on disks

Unit assignment for $k=5$, $m=2$

resource	data					parities	
	r_0	r_1	r_2	r_3	r_4	r_5	r_6
units	0	3	6	9	12	15	18
	1	4	7	10	13	16	19
	2	5	8	11	14	17	20

number of units per resource (ω)

- $\omega = 3$,
- generally $2^\omega > k + m$



Cauchy Reed/Solomon: Equation-based definition

Coding algorithm is an execution of several equations:

- either: instant application on data
- or: store and transform equations, apply them on a sequence of code words

Example of a 5+2 Reed/Solomon code:

direct encoding (45 XOR op.)

```
15 = XOR(2, 3, 4, 5, 7, 9, 11, 12)
16 = XOR(0, 2, 3, 7, 8, 9, 10, 11, 13)
17 = XOR(1, 3, 4, 6, 8, 10, 11, 14)
18 = XOR(0, 2, 4, 6, 7, 8, 11, 12, 13)
19 = XOR(0, 1, 2, 4, 5, 6, 9, 11, 14)
20 = XOR(1, 2, 3, 5, 6, 7, 10, 12)
```

direct decoding (42 XOR op.)

```
6 = XOR(0, 5, 17, 18, 20, 13, 14)
7 = XOR(1, 3, 5, 15, 17, 18, 19, 20, 12, 13)
8 = XOR(2, 4, 16, 19, 20, 12, 13, 14)
9 = XOR(2, 3, 4, 15, 16, 17, 19, 12, 13)
10 = XOR(0, 2, 5, 15, 19, 20, 14)
11 = XOR(1, 3, 15, 16, 18, 20, 12)
```

iterative encoding (33 XOR op.)

```
15 = XOR(B, C, D)
16 = XOR(D, E, F)
17 = XOR(3, 4, 8, E, H)
18 = XOR(2, 4, 6, 7, C, F)
19 = XOR(0, 2, 9, 11, B, H)
20 = XOR(5, 7, 10, 12, A, G)
A = XOR(2, 3)      E = XOR(10, 11)
B = XOR(4, 5)      F = XOR(0, 8, 13)
C = XOR(11, 12)    G = XOR(1, 6)
D = XOR(7, 9, A)   H = XOR(14, G)
```

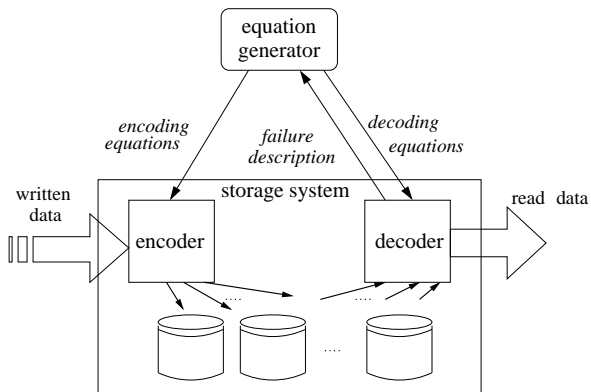
iterative decoding (29 XOR op.)

```
6 = XOR(B, C)
7 = XOR(5, C, D, F)
8 = XOR(19, 14, A, G)
9 = XOR(3, 17, 13, D, G)
10 = XOR(2, 20, B, D)
11 = XOR(15, 16, 18, 20, F)
A = XOR(20, 13)    D = XOR(15, 19)
B = XOR(0, 5, 14)  F = XOR(1, 3, 12)
C = XOR(17, 18, A) G = XOR(2, 4, 12, 16)
```

Equations for coding

Separation:

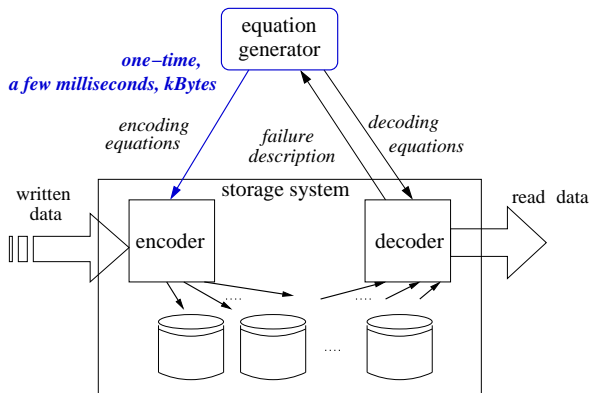
- equation preparation
- equation interpretation for coding



Equations for coding

Separation:

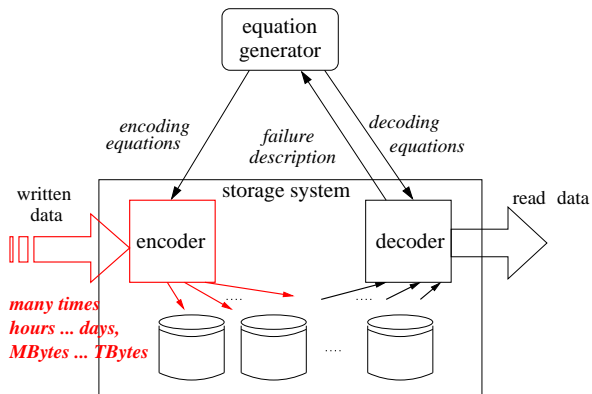
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Equations for coding

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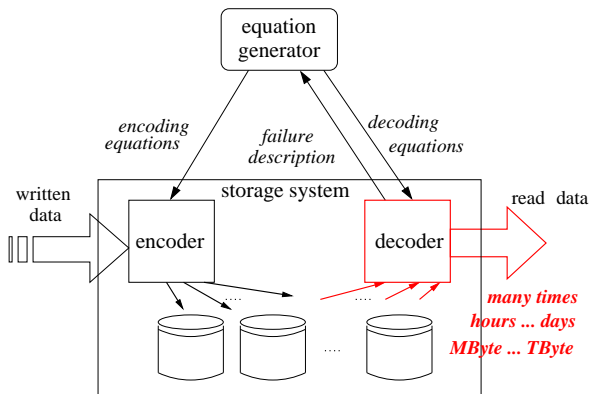
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Equations for coding

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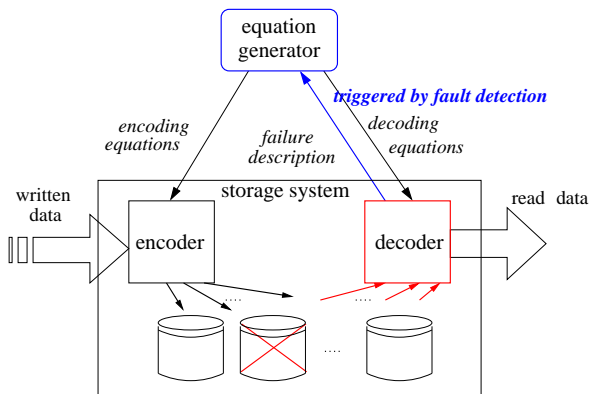
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Equations for coding

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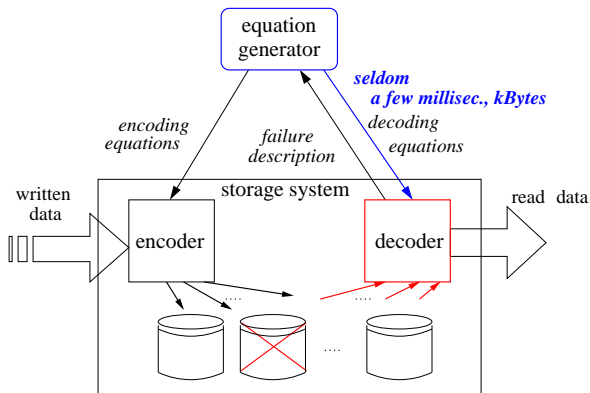
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Equations for coding

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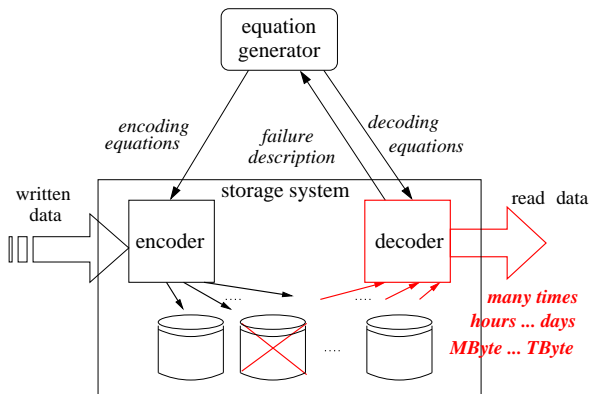
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Equations for coding

Separation:

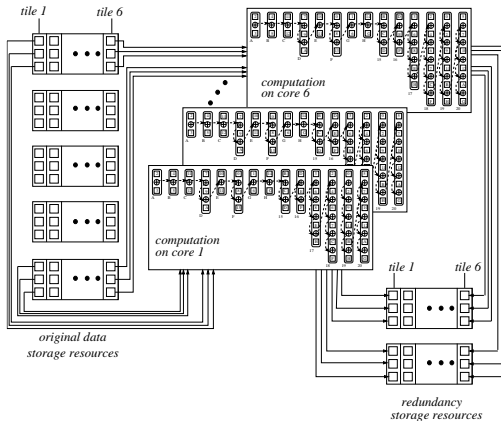
- equation preparation
- equation interpretation for coding



Parallel Coding

Obvious parallelism: block parallel coding

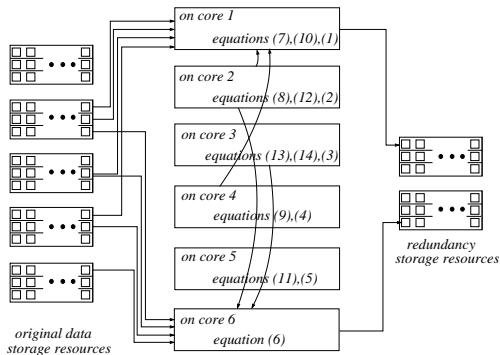
- same coding function on different data blocks
- a core interprets all equations
- a core streams only a part of the input data



Parallel Coding - Equation oriented

Equation-oriented coding

- a core interprets dedicated equations
- a core streams data which is referred by the dedicated equations



Parallel Coding Schedules

encoding and decoding equations extended to schedules

schedule:

- equations assigned to cores
- XOR ops assigned to time steps
- data dependencies resolved

cores	steps					
	1	2	3	4	5	6
1		7 ⊕ 9	A	B ⊕ C ⊕ D		
	2 ⊕ 3					
2		0 ⊕ 8 ⊕ 13		F ⊕ E ⊕ D		
	4 ⊕ 5					
3		14 ⊕ G	3 ⊕ 4 ⊕ 8 ⊕ E ⊕ H			
	1 ⊕ 6					
4		2 ⊕ 4 ⊕ 6 ⊕ 7 ⊕ C ⊕ F				
	10 ⊕ 11					
5		0 ⊕ 2 ⊕ 9 ⊕ 11 ⊕ B ⊕ H				
	11 ⊕ 12					
6		7 ⊕ 5 ⊕ 10 ⊕ 12 ⊕ A ⊕ G				

Schedule preparation: stacking of equations

Parallel Coding Schedules

Stacking of equations

Encoding equations,
33 XOR operations

- Terminal equations

$$15 = \text{XOR}(B, C, D)$$

$$16 = \text{XOR}(D, E, F)$$

$$17 = \text{XOR}(3, 4, 8, E, H)$$

$$18 = \text{XOR}(2, 4, 6, 7, C, F)$$

$$19 = \text{XOR}(0, 2, 9, 11, B, H)$$

$$20 = \text{XOR}(5, 7, 10, 12, A, G)$$

- Temporary equations

$$A = \text{XOR}(2, 3)$$

$$B = \text{XOR}(4, 5)$$

$$C = \text{XOR}(11, 12)$$

$$D = \text{XOR}(7, 9, A)$$

$$E = \text{XOR}(10, 11)$$

$$F = \text{XOR}(0, 8, 13)$$

$$G = \text{XOR}(1, 6)$$

$$H = \text{XOR}(14, G)$$

cores	steps							
	1	2	3	4	5	6		
1		7	⊕ 9	⊕ A	B	⊕ C	⊕ D	
2	2	⊕ 3						
3				F	⊕ E	⊕ D		
4	4	⊕ 5						
5			3	⊕ 4	⊕ 8	⊕ E	⊕ H	
6	1	⊕ 6		14	⊕ G			
7								
8	2	⊕ 4	⊕ 6	⊕ 7	⊕ C	⊕ F		
9	10	⊕ 11						
10			0	⊕ 2	⊕ 9	⊕ 11	⊕ B	⊕ H
11	11	⊕ 12						
12								
13	7	⊕ 5	⊕ 10	⊕ 12	⊕ A	⊕ G		
temporary units required								
G A B,A D,G H								
F,C E								
temporary units available								
A,B, H D,F								
C,G,E								

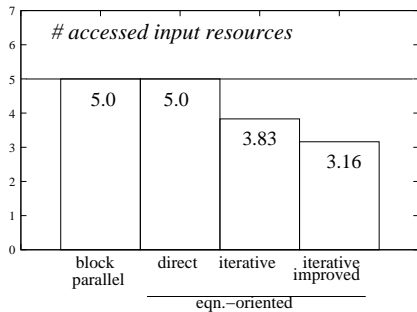
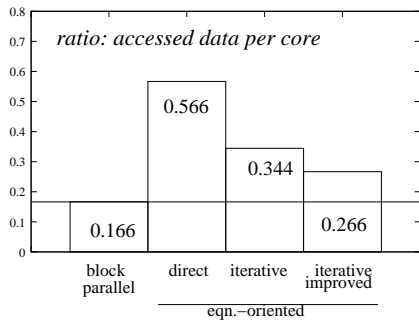
Evaluation

Question: Is equation-oriented parallel coding beneficial?

Criteria:

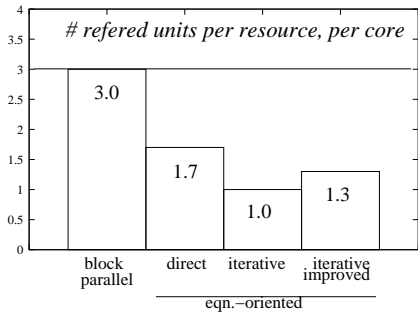
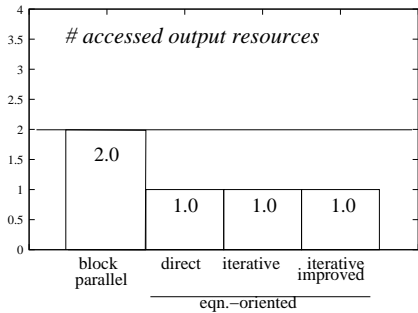
- accessed data per core
- number of referenced storage resources (input, output)
- number of referenced units per resource and per core
- multiplicity of references
- number of temporary results taken from other cores
- number of time steps of a schedule
(under absence of access delays, and synchrony of XOR operations)

Evaluation



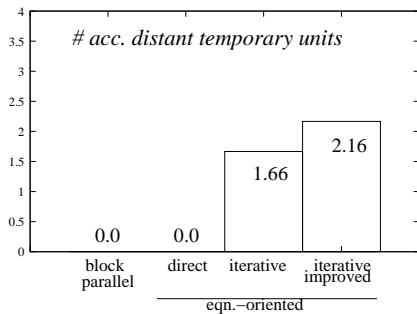
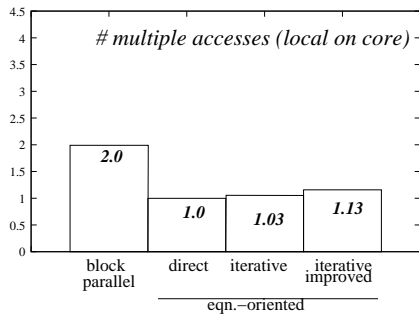
Lower values are better

Evaluation



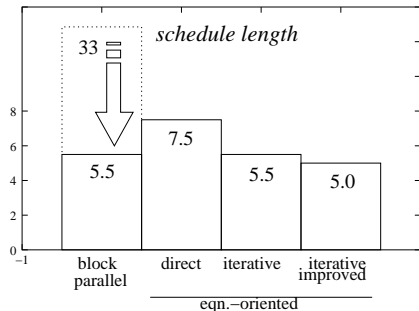
Lower values are better

Evaluation



Lower values are better

Evaluation



Equation-oriented parallel coding:

- iterative equations only!
- improve: selecting good Cauchy matrices

Performance benefits:

- minimal schedule length
- multiple accesses reduced
- locality of accesses (resources, units)

Performance obstacles:

- access to temporary results from other cores

Summary

- Cauchy-Reed/Solomon code: XOR based
- Decomposition of coding into several parts, described by equations
- Equations: parameterize the encoding and decoding function
- Schedules: pre-calculated placement of equations on cores
- Iterative schedules:
concentration of data accesses of a core on local regions
- Advantage for software-based coding performance on multicore processors