



On Configuring Hierarchical Storage Structures

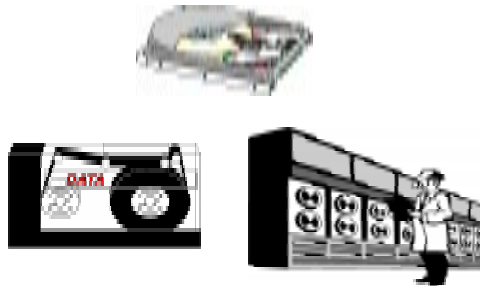
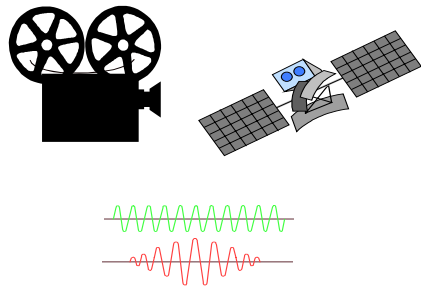
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Configuration Problem



Growing Data Sets

Numerous Storage Alternatives

User Requirements

Challenges & Contributions

■ Cost effective configuration (i.e., balanced-design):

- Configuration planer to find cost effective HSS designs for a range of throughputs, DB sizes, and access distributions

■ Continuous display using HSS's:

- Display Caching (DC) technique to support continuous display from tape devices
- Object Caching (OC) of hot objects at lower memory levels (e.g., on disks)

Configuration Categories



	DRAM	Disk	Tape
D-based	DC	Complete DB	-
M-based	Complete DB*	-	-
T-based	DC	DC	Complete DB
MD-Hierarchy	DC & OC	Complete DB	-
MT-Hierarchy	DC & OC	DC	Complete DB
DT-Hierarchy	DC	DC & OC	Complete DB
MDT-Hierarchy	DC & OC	DC & OC	Complete DB

M: Memory

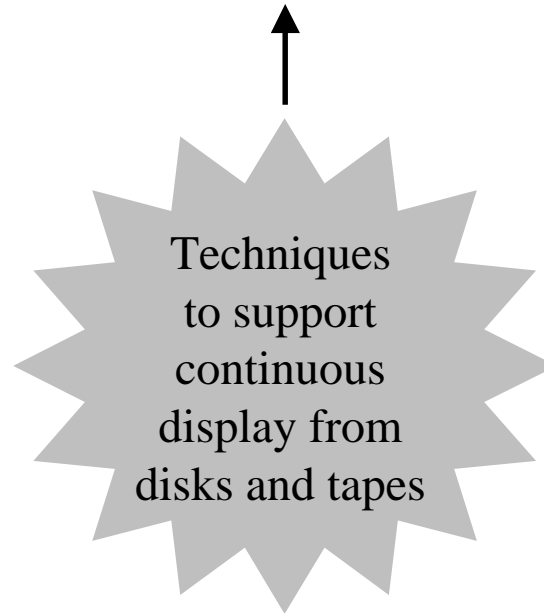
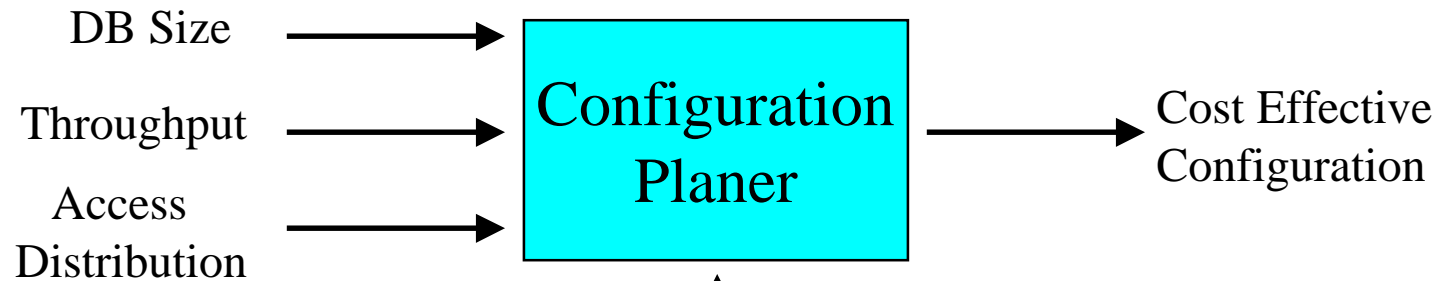
D: Magnetic Disk

T: Magnetic Tape

DC: Display-Caching

OC: Object-Caching

Configuration Planer



Display Caching



- Display Caching (DC) technique is used to bridge the gap between the retrieval behavior of the tape juke box (D_T) and the display requirement of objects (C):
 - How much data should be cached?
 - Where the data is cached?
 - When the display should start?
- Production Consumption Ratio ($PCR = D_T / C$)

Display Caching (2): PCR < 1

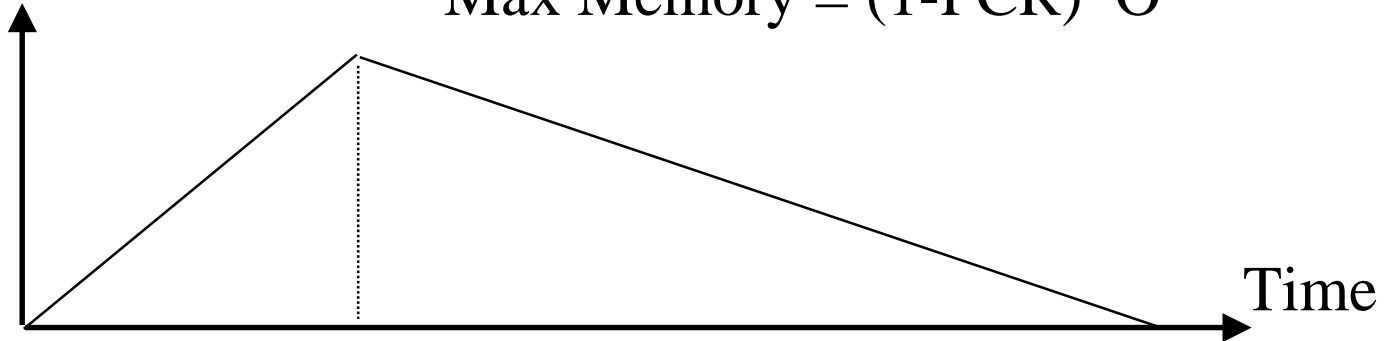


Object Retrieval

Object Display

Memory
Requirements

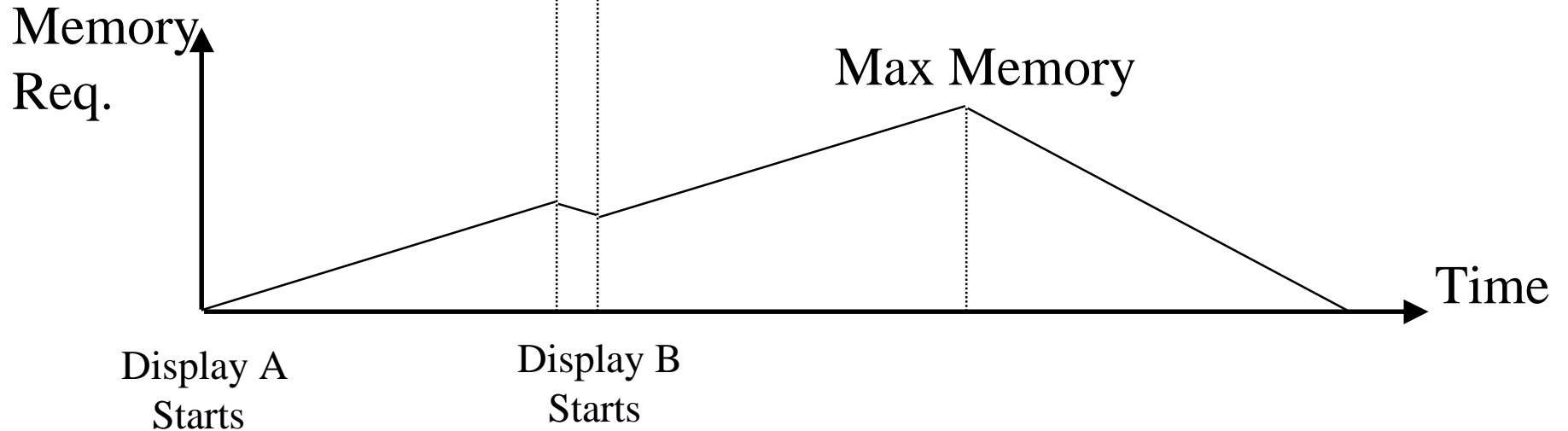
$$\text{Max Memory} = (1 - \text{PCR}) * O$$



Display

$$\text{Starts } (t_0 = (1 - \text{PCR}) * O / DT + \delta)$$

Display Caching: $PCR > 1$



Display Caching: PCR > 1



- Number of simultaneous displays from the display cache is:

- $N_T = \text{MAX}(1, \text{floor}(O / (T_{\text{Access}} + O/D_T)))$

- Maximum memory requirements:

- Max Memory =
 $\sum_{k=0}^{N_T-1} ((1 - 1/\text{PCR})^k * O - k * (T_{\text{Access}} + O/D_T) * C)$, for all
 $k=0 \rightarrow N_T-1$

Results



	100 Gbytes	500 Gbytes	1 Tbytes	5 Tbytes	10 Tbytes
10	M=1 D=3 T=1	M=1 D=3 T=1	M=1 D=4 T=2	M=1 D=5 T=6	M=1 D=4 T=11
100	M=4 D=11	M=5 D=38 T=1	M=5 D=44 T=2	M=3 D=20 T=6	M=6 D=30 T=11
1000	M=33 D=36	M=34 D=56	M=39 D=111	M=45 D=420 T=6	M=44 D=198 T=63
5000	M=163 D=179	M=163 D=179	M=163 D=179	M=174 D=556	M=190 D=1022 T=11
10000	M=1473 D=11	M=325 D=358	M=325 D=358	M=331 D=556	M=348 D=1111

M,D, & T represent number of modules